Personalized Feedback-based Customer Incentives in Automated Demand Response

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Introduction

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- Demand Response (DR) programs for curtailing energy consumption in critical times for the grid are becoming common
- Automated DR (ADR) automates the response process of the customer to the DR signals by means of electric controls installed at the customer premises
- ADR rebates are defined mostly statically and based on
 - either the costs of ADR equipment or the cost per unit of energy at peak times
- Two problems with that:
 - First, the utility loss for the customer due to curtailed energy may include other aspects, such as actual needs, sensitivity to personal-comfort loss, etc.
 - Endowment may fall short for engagement
 - Second, ADR programs currently do not take into account the customer satisfaction from the provided endowment for load curtailment
 - Unsatisfied customers may not renew contracts

Our Objectives

- 1) Find ADR endowments that satisfy customers for a specific load curtailment
 - For non-purely rational customers
 - Even when user utility functions are not known
- 2) Keep incentive cost as low as possible or within a specific budget
 - Trade-off between load curtailment, incentive cost and customer satisfaction

System Model

System Model

- A customer *i* enjoys net benefit U_i (i.e., user satisfaction minus energy cost) from consuming baseline energy q_i^0
- An energy-consumption reduction ΔQ_i in specific time periods according to an ADR contract results to a net-benefit loss $\Delta U_i = -\eta_i U_i$
 - Different per customer *i*
- In return, the customer *i* receives an endowment *b_i* by the utility company

User Utility Model

• Generic user utility model: $u_i = g(U_i, U_{-i})$

> NB of others

net benefit (NB) from power consumption of customer i

- User utility difference:
 - $\Delta u_i = \hat{g}(\Delta U_i, \Delta U_{-i}) + b_i$
- Specific instance of user utility model: *altruism u_i* = (1 *y_i*) *U_i* + *y_iU_{-i}*
 - $-\gamma_i \in [0, 1]$ is the degree of altruism
- Then: $\Delta u_i = -(1 \gamma_i)\eta_i U_i + \gamma_i \Delta U_{-i} + b_i$

DR Designer's Problem

Optimization goals

- I. Maximize customer satisfaction α for a specific net benefit reduction η_i due to load curtailment for each customer *i* within a budget limit *B* for endowments
- II. Minimize total endowment cost for a lower-bound η in the net-benefit loss due to load curtailment of each customer and for a lower-bound α in customer satisfaction

Full-info solutions: Uniform Endowment

- Full-information on user utilities
- Observe that customer satisfaction ratio α is monotonic in the uniform endowment b
- Problem (I) can be solved by sorting all consumers with respect to Δu_i of each customer *i* after using maximum endowment b = B/N and count how many of them are positive
- Problem (II) can be solved by sorting all consumers with respect to Δu in a descending order for b=0 and then solve the equation $\Delta u = 0$ at position k to find b, with $k = \alpha + N$. This is the value of b that makes α % customers having $\Delta u \ge 0$

 Δu_{s} Δu_2 ∆u₃ Δu_{10} Solve $\Delta u_4 = 0$ Δuα Δu_7 ∆u₁ $\alpha = 80\%$ ∆u₄ satisfied ∆u₆ ∆u₅ 10

N=10

Full-info solutions: Personalized Endowment

- Full-information on user utilities
- If personalized incentive b_i per customer i is employed, then problem (I) is again solved as described above, while problem (II) is solved as follows:
 - For each customer *i*, calculate the personalized incentive that renders $\Delta u_i = 0$
 - Sort the list of customers based on their personalized incentive in ascending order
 - The minimum total incentive required for satisfying $\alpha \cdot 100\%$ customers is given by summing the top-(αN) personalized endowments

Hidden Info \rightarrow Customer Feedback

- Customers provide feedback on satisfaction
 - In a ballot



– Personalized



- It can also be strategic!



Distributed Algorithm

- At each round *t*, the DR designer sets a b_t and each customer *i* responds to it with feedback $v_{\dot{r}t+1}$, which collectively result to a mean satisfaction level α_{t+1} for the received incentive at the next round
- The feedback $v_{i,t+1}$ of customer *i* at round t + 1 is determined by the sign of:

$$\Delta u_{i,t+1} = \hat{g}(\Delta U_i, \Delta U_{-i,t}) + b_t$$

 Employing gradient ascent, the DR designer selects b_{t+1} for the round t+1 as follows:

$$b_{t+1} = \begin{cases} \max\{b_t + \Delta \alpha \cdot \kappa, 0\} \\ b_t + \kappa, \end{cases}$$

Stopping Criteria

- Problem (I): If $\Delta \alpha / \alpha_t < \Delta b / b_t$ or $b_t \ge B$, then stop iterations
- Problem(II): if $\alpha_t \ge \underline{\alpha}$, then stop iterations

Estimating $\widetilde{\Delta U}_{-i,t}$

- Assuming $\frac{\sum_{i=1}^{N} (1 \gamma_i) \eta_i U_i}{N} \approx (1 \bar{\gamma}) \overline{\Delta U}$
- We obtain $\overline{\Delta u}_{t+1} = \overline{\Delta U}_t + b_t$
- Observe that $\overline{\Delta u}_{t+1}$ and $\eta_{max}(a_{t+1}-1)+b_t$ have the same output sets and move similarly according to $\Delta u_{i,t+1}$ values
- Then, approximate that $\overline{\Delta u}_{t+1} = \eta_{max}(a_{t+1} 1) + b_t$
- It follows that

$$\overline{\Delta U}_t \approx \eta_{max}(a_{t+1} - 1)$$

Strategic Feedback

- However, customers have incentive to lie on their satisfaction
- DR mitigation policy:
 - The DR designer b sets an upper bound on the budget B for endowments that is unknown to the customers
 - If b_t becomes infeasible, then no endowment is provided (load curtailment is still sustained)
- Then, user utility difference function for customer *i* becomes

$$\Delta u_{i,t+1} = -(1 - \gamma_i)\eta_i U_i + \gamma_i \overline{\Delta U}_{-i,t} + \Pr[b_t < B | z_i \text{ lies}]b_t$$

 $\eta_{max}(a_{t+1}-1)$

of times that customer lied

Customer Targeting

Customer Targeting

- Recall that same energy consumption reduction ΔQ results to a different net benefit loss fraction η_i for each customer *i*
 - According to internal individual function $h_i(\Delta Q_i)$ of each customer *i*
- Then, customer utility difference is given by

$$\Delta u_i = -(1 - \gamma_i)h_i(\Delta Q_i)U_i + \gamma_i\overline{\Delta U}_{-1} + b_i(\Delta Q_i)$$

• Assume discrete levels of consumption reduction in $H = \left\{ \frac{1}{N} \Delta Q, \frac{2}{N} \Delta Q, \dots, \theta_{max} \Delta Q \right\}$

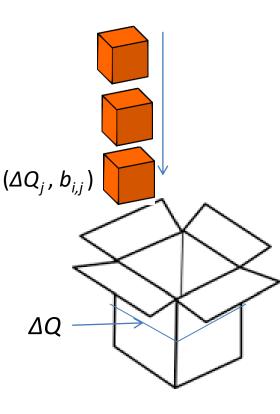
Finding Endowments for Targeting

- Problem: For each $\Delta Q_j \in H$, find $b_{i,j}$
- **Full info:** Simply solve $\Delta U_i(b) \ge 0$ for each customer *i*

- Hidden info, yet customer feedback individually observable or not :
 - Assume each ΔQ_j a uniform reduction for all customers and employ the distributed algorithm for determining either personalized or uniform $b_{i,j}$ for each customer *i*
 - Entails approximation due to altruism

Targeting Algorithm

- Given $(\Delta Q_j \text{ KWh}, b_{i,j} \in)$ pairs in list L
- Sort them based on $\Delta Q_j / b_{i,j}$ in <u>decreasing</u> <u>order</u>
- Add pairs from list *L* into a list *S* until next item exceeds desired total ΔQ
- Then, if desired total ∆Q has been reached in S, you are done
- Otherwise, from remaining items in L find the one that completes exactly ΔQ, if exists, and add to S; if not, add to S the cheapest item, so that desired ΔQ is overfilled



List L



It can find optimal solution!

"When it is possible our targeting algorithm to fill the bag with exactly ΔQ , it finds an optimal solution."

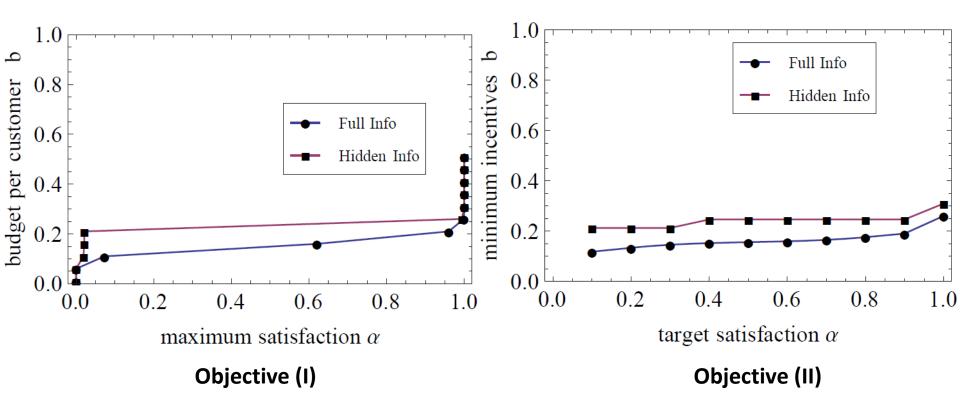
- Sketch of Proof:
 - By contradiction: trying to replace one of the items in bag *S*, as selected by the algorithm so that ΔQ load is curtailed, with one or multiple other ones results in higher total incentive cost

Evaluation

Evaluation Setup

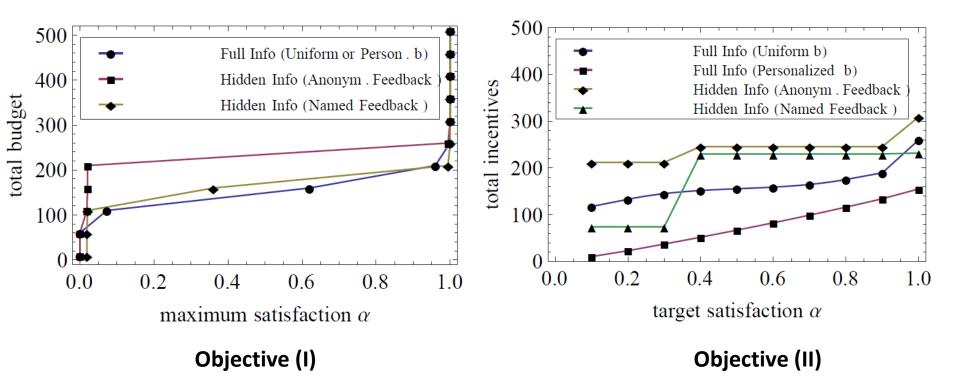
- *N* = 1000 customers
- Altruism for customer *i*: $\gamma_i \sim U(0, 1)$
- Net benefit loss of customer *i*: $\eta_i \sim U(0.1, 0.3)$, unless otherwise specified
- The DR designer is assumed to have guessed semi-correctly $\eta_{max} = 0.5$
- Satisfaction of customer *i* for her nominal energy consumption: U_i
 ~ N(0.8, 0.1)
- *U_i* assumed normalized by maximum net benefit, so is *b*
- Nominal consumption q_i^0 is 1 for all customers

Uniform Load Reduction: Anonymous Feedback



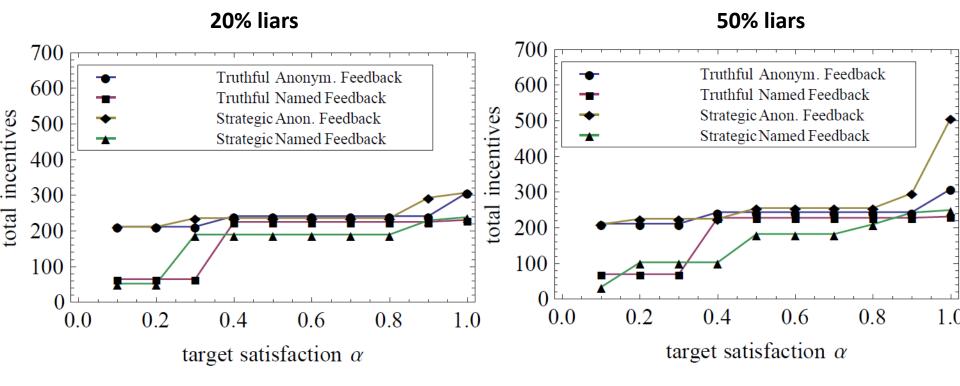
- In case of hidden info, the distributed algorithm finds uniform endowment close to those of the full-info case
- Thus, approximation of $\Delta U_{-i,t}$ is good enough

Uniform Load Reduction: Named Feedback

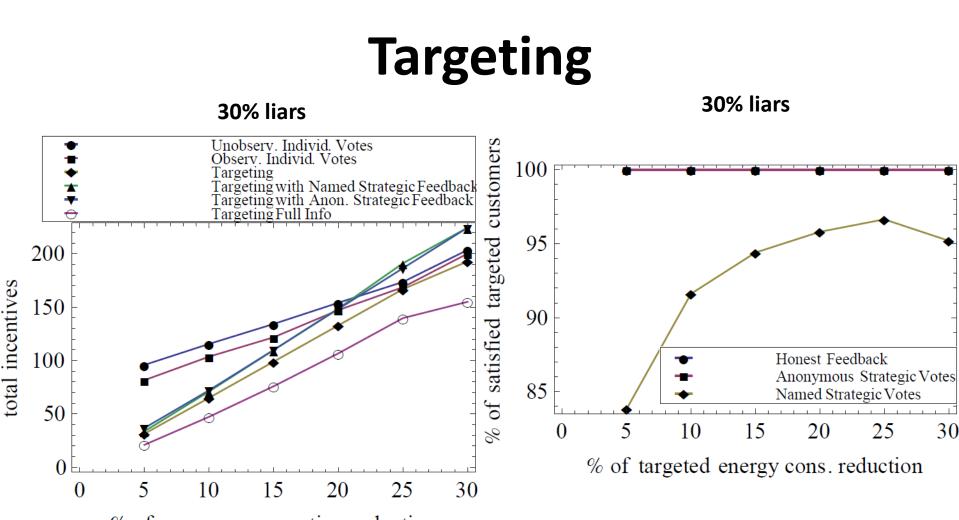


Named feedback almost reveals hidden info for objective (I), while it moderately improves endowment cost for objective (II)

Strategic Lying



Strategic lying is successfully mitigated



% of energy consumption reduction

- Targeting lowers significantly endowment costs in case of hidden info for low fractions of load curtailment
- Not affected significantly by strategic lying
- Targeted customers are kept satisfied, even in presence of lying!

Conclusions

Conclusions

- We proposed algorithms for calculating satisfactory ADR endowments for uniform or personalized energy-load reduction for non-rational customers in the cases of both full and hidden info on user utilities
 - In case of hidden info on user utilities, we employed anonymous and named feedback on customer satisfaction, which may be strategic or not
- Our evaluation has shown the effectiveness of the various algorithms for all cases
 - even in the presence of high fractions of strategic liars among customers
- Customer targeting is preferable for low (<20%) desired energy consumption reductions, even for hidden information on user utility functions and even in the presence of 30% strategic liars
- Our formulation and approach are generic-enough to consider different user utility functions
- As a future work, we plan to consider different behavioral factors in the user utility

Thank you for your attention

Any questions?

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http://stecon.cs.aueb.gr/research/energy-environment/





Support Slides

Estimating $\widetilde{\Delta U}_{-i,t}$

- It holds $-1 \le -U_i \le 0$
- Since $\eta_i \in [0, 1]$, it follows that

 $-\eta_i \leq \Delta U_i \leq 0$ and $-\eta_{max} \leq \Delta U \leq 0$,

• Therefore and since $\gamma_i \in [0, 1]$, it is true that

$$b_t - \eta_{max} \le \Delta u_{i,t+1} \le b_t$$
, $\forall i \in N$

- Recall that $\overline{\Delta U}_{-i} = -\frac{\sum_{\substack{j=1\\j\neq i}}^{N} \eta_j U_j}{N}$
- Adding down utility differences and dividing by *N*, we have

$$\overline{\Delta u}_{t+1} = -\frac{\sum_{i=1}^{N} (1 - \gamma_i) \eta_i U_i}{N} + \bar{\gamma} \overline{\Delta U}_t$$